

# The Perfect Functional Programming Course

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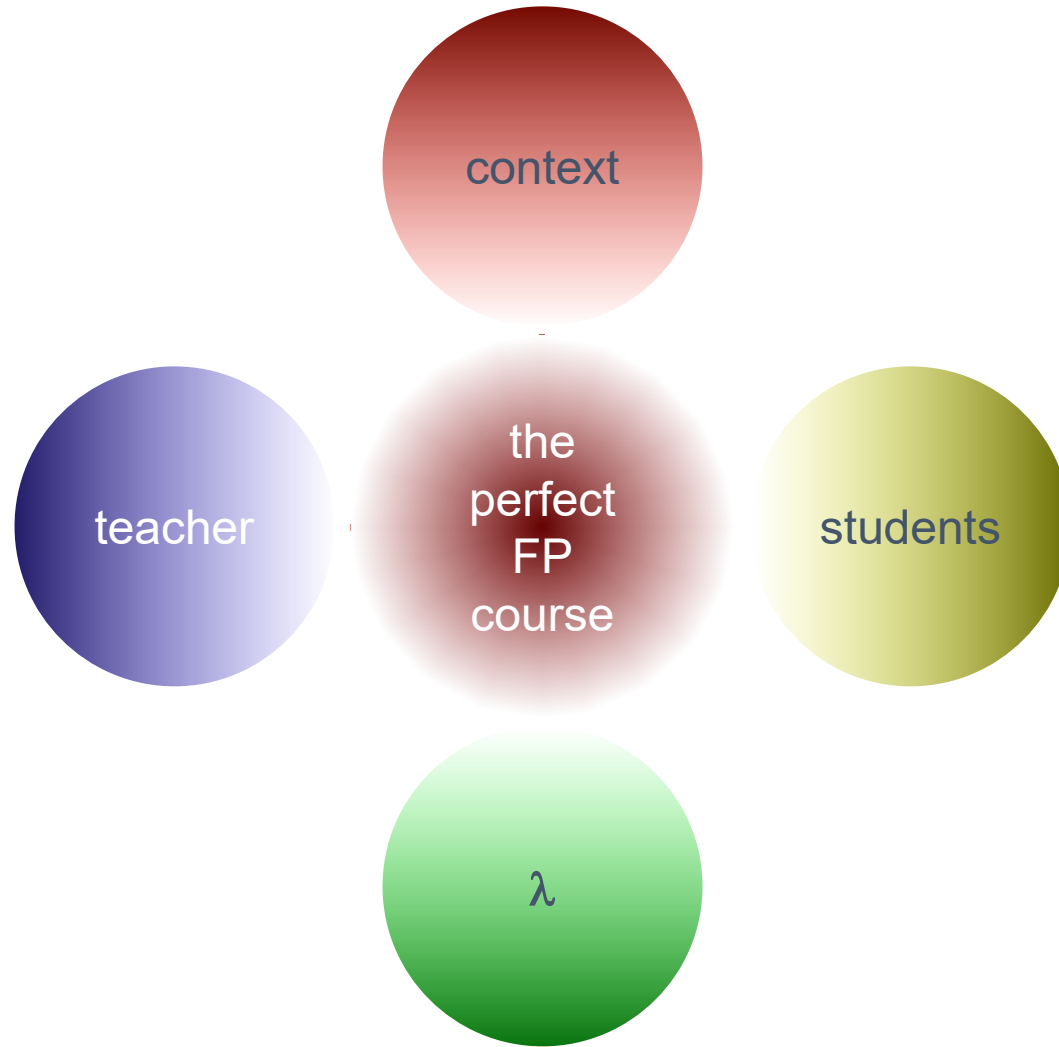
# What this talk is about

- Share my experience in teaching functional programming
- A perfect (functional programming) course?
- Effective remedies to deal with hurdles that students face in functional programming
- Discuss engaging student project ideas
- Wrap up

# Who am I

- Computing science education at Radboud University (1985 – 1991)
- PhD Thesis ‘Interactive Functional Programs’ (1996)
- Lecturer at Radboud University (2000 – )
- Teacher of programming language related courses:
  - imperative programming (bachelor 1<sup>st</sup> year introductory course, C++)
  - functional programming (bachelor 1<sup>st</sup> year, 2<sup>nd</sup> year versions, also for AI students, Clean)
  - compiler construction (master course, language agnostic)
  - advanced programming (master course, Clean, SaC)

# The big picture



# Context



- Placement in curriculum
- Course end terms
- Resources
  - study credits
  - study activities
  - educational forms
  - teaching assistants
  - programming language (support)
- Students
  - background
  - required knowledge
  - problem solving skill set

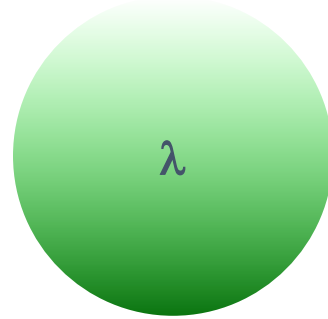
# Students



students

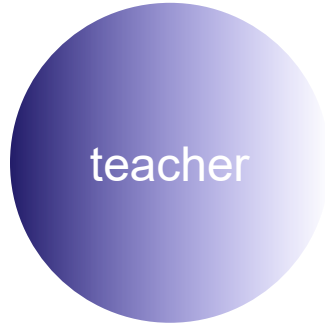
- “What’s in it for me?”
  - extend problem solving skill set
  - revisit known problems and solve them again
  - visit new problems and solve them
  - make reasoning about programs tractable
- Opportunities
  - relate to other courses
  - ‘make the connection’

$\lambda$

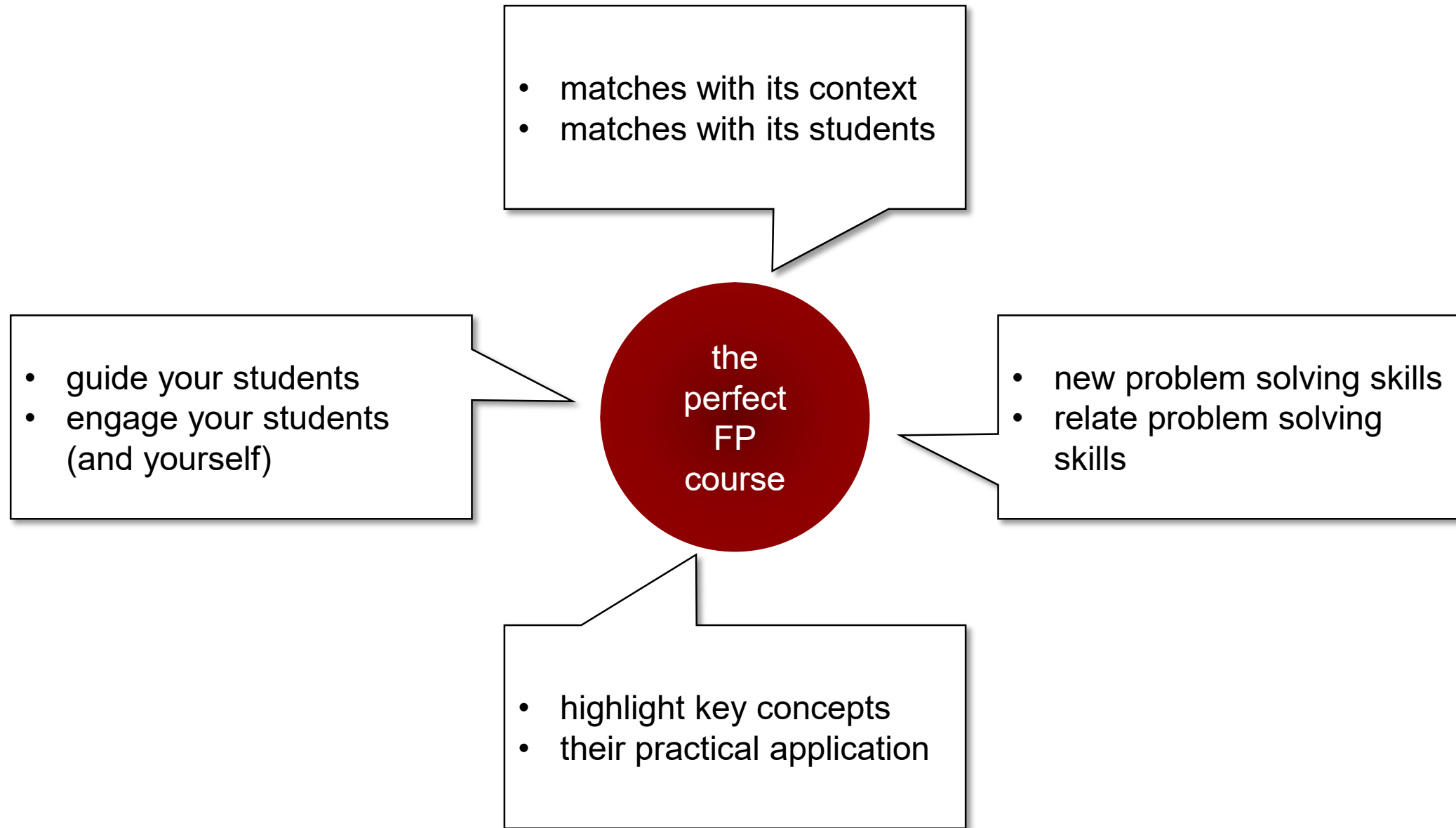


- “What’s in it for me?”
- Liberate programmers from the von Neumann style
- Show the elegance of programming with equations
  - expressive data types and type systems
  - pattern matching / case distinctions
  - higher-order functions
  - evaluation strategies
  - (equational) reasoning instead of debugging

# Teacher



- “What’s in it for me?”
- Show students that only a few key concepts lead to a rich problem solving tool set
- Less exam grading time than with imperative / object oriented programming languages
- Engage your students and yourself
  - Why Functional Programming Matters [Hughes <sup>1984</sup>]
  - Introduction to Functional Programming [Bird, Wadler <sup>1988</sup>]
  - Purely Functional Data Structures [Okasaki <sup>1998</sup>]
  - Functional Programs as Executable Specifications [Koopman <sup>1999</sup>]
  - Haskell School of Expression [Hudak <sup>2000</sup>]
  - How to Design Programs [Felleisen et al <sup>2001</sup>]
  - SoccerFun [Achten <sup>2008</sup>]
  - Animated Problem Solving [Morazán <sup>2022</sup>]



# A perfect functional programming course?

# A perfect functional programming course?

## ① Rewriting

### Why:

- get into the basics of functional programming
- help students understand the mechanics
- seed for higher-order functions
- seed for evaluation strategies
- seed for equational reasoning

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$$\begin{aligned} & 8 \div 2(2 + 2) \\ &= (8 \div 2) \bullet (2 + 2) \\ &= \underline{(8 \div 2)} \bullet (2 + 2) \\ &= 4 \bullet \underline{(2 + 2)} \\ &= \underline{4 \bullet 4} \\ &= 16 \end{aligned}$$

# A perfect functional programming course?

- ① Rewriting
- ② Type inference

## Why:

- remove the mysticism around types
- help students understand type error messages
- polymorphic functions ‘for free’

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types

## Why:

- intuitive way to model problem domains
- pattern matching on custom types
- recursive problem domains
- overloaded functions

# Modelling a problem domain: ancestry

- We want to model the following facts regarding a person's ancestry:
  1. a person has a name
  2. a person has a date of birth
  3. a person has a blood group
  4. a person has two biological parents
  5. a date consists of a year, month, day
  6. a biological parent is either a **person** or is unknown

ancestry is recursive

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types
- ④ Lists and list comprehensions

## Why:

- workhorse of functional programming
- pattern matching on lists
- lists and recursion
- list comprehensions
- seed for higher-order functions

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types
- ④ Lists and list comprehensions
- ⑤ Higher-order functions

## Why:

- key concept of functional programming
- key to abstraction and reusable code
- students get used to having functions as first-class citizens

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types
- ④ Lists and list comprehensions
- ⑤ Higher-order functions
- ⑥ Evaluation strategies

## Why:

- deepen understanding of rewriting
- discuss and compare evaluation strategies
- referential transparency
- seed for equational reasoning

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types
- ④ Lists and list comprehensions
- ⑤ Higher-order functions
- ⑥ Evaluation strategies
- ⑦ (Equational) reasoning

## Why:

- software correctness matters
- deepen abstraction
- deepen reusable code
- advocate reasoning over debugging

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types
- ④ Lists and list comprehensions
- ⑤ Higher-order functions
- ⑥ Evaluation strategies
- ⑦ (Equational) reasoning
- ⑧ Case studies or projects

## Why:

- engage students
- not necessarily at the end

# A perfect functional programming course?

- ① Rewriting
- ② Type inference
- ③ Custom data types
- ④ Lists and list comprehensions
- ⑤ Higher-order functions
- ⑥ Evaluation strategies
- ⑦ (Equational) reasoning
- ⑧ Case studies or projects

## Room for deepening:

- context, students, language, teacher
- programming pattern abstractions:
  - monads, functors, monoids, ...
  - type constructor classes, gadt's, ...
- programming pragmatics:
  - interactive programs (console, GUI, graphics)
  - file / data processing
  - evaluation strategies and side-effects
- algorithms and data structures:
  - (binary) search trees, AVL trees, ...
- compiler / interpreter / DSL:
  - parser combinators, grammars, interpreters
- ...

# Common hurdles for students in functional programming

# Common hurdles for students in functional programming

- Types and type error messages
- List processing functions
- Get a grip on higher-order functions

# Types and type error messages

- Students struggle with typing functions and type error messages
- Effective remedy is to explain typing as a logical puzzle

# Type inference is a logical puzzle

- Rules of the game (Clean book, section 1.5.7):
  - everything has one type
  - in a function the same argument has the same type in all function alternatives
  - all alternatives of a function have the same type
- Strategy to solve the puzzle:
  1. give each argument and result a fresh (lowercase) name-**?** pair and add them in a table
  2. check all pattern matches to extract type structure information
    - replace **?** with the found information and create fresh name-**?** pairs if necessary
    - replace in the function type and the table every occurrence of that name with the found information
  3. check all<sup>1</sup> sub-expressions to extract type information
    - replace **?** with the found information and create fresh name-**?** pairs if necessary
    - replace in the function type and the table every occurrence of that name with the found information
- Any remaining name-**?** pair represents a polymorphic type variable

---

<sup>1</sup> compiler has to check all sub-expressions to detect possible errors

# Type inference is a logical puzzle

```

module MySecondCleanProgram

import StdEnv

1 space 0 = ""
  space n = " " ++ space (n-1)

2 pyramid` n "" = ""
  pyramid` n text = space n ++ text ++ "\n" ++
                    pyramid` (n+1) (text%(1,size text-2))

```

```

(-)  :: Int Int -> Int
(+)  :: Int Int -> Int
(+++) :: String String -> String
(%)  :: String (Int,Int) -> String
size :: String -> Int

```

**step 1:** fresh names for the arguments and result

# Type inference is a logical puzzle

```

module MySecondCleanProgram

import StdEnv

space :: a -> b
space 0 = ""
space n = " " ++ space (n-1)

pyramid` n "" = ""
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```

1

2

```

(-)  :: Int Int -> Int
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(+++) :: String String -> String
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size :: String -> Int

```

```

a = Int
b = ?

```

**step 2:** check all pattern matches

# Type inference is a logical puzzle

```

module MySecondCleanProgram

import StdEnv

space :: Int -> b
space 0 = ""
space n = "" +++ space (n-1)

pyramid` n "" = ""
pyramid` n text = space n +++ text +++ "\n" +++
                  pyramid` (n+1) (text%(1,size text-2))

```

1

2

```

(-)  :: Int Int -> Int
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(+++) :: String String -> String
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```

```

a = Int
b = String

```

**step 3:** check all function alternatives

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```

2

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pyramid` n "" = ""
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```

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a = ?
b = String
c = ?

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```

a = ?
b = String
c = String

```

**step 3:** check all function alternatives

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a = Int
b = String
c = String

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**step 3:** check all function alternatives



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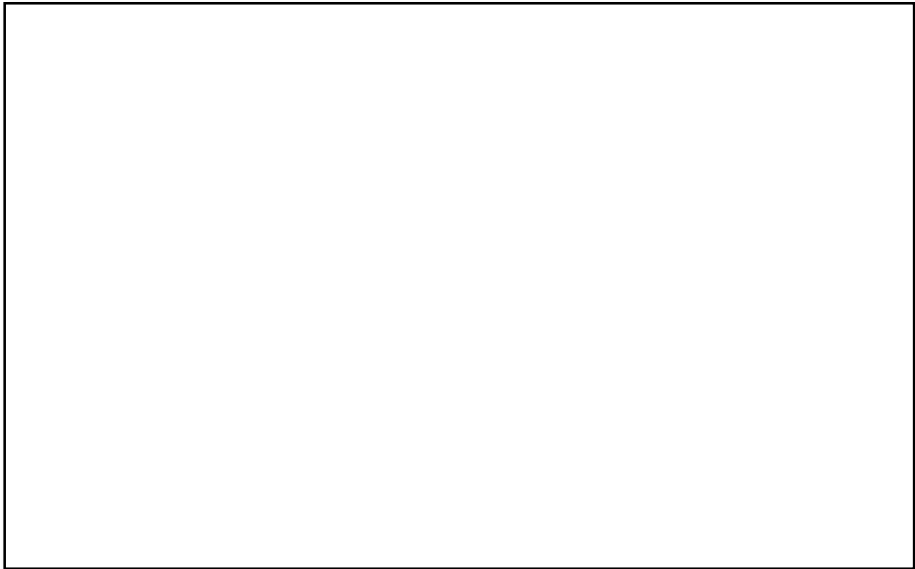
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# Type inference with tuples

```
implementation module MyTupleFunctions
```

- 1 `first (x,y) = x`
- 2 `swap (x,y) = (y,x)`

**step 1:** fresh names for the arguments and result

# Type inference with tuples

```
implementation module MyTupleFunctions
```

- ① `first :: a -> b`  
`first (x,y) = x`
- ② `swap (x,y) = (y,x)`

```
a = ?  
b = ?
```

**step 1:** fresh names for the arguments and result

# Type inference with tuples

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implementation module MyTupleFunctions
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① `first :: a -> b`  
`first (x,y) = x`

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a = ?

b = ?

**step 2:** check all pattern matches

# Type inference with tuples

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implementation module MyTupleFunctions
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```
first :: a -> b
```

① 

```
first (x,y) = x
```

② 

```
swap (x,y) = (y,x)
```

```
a = (c, d)
```

```
b = ?
```

```
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```
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a = (c, d)
```

```
b = c
```

```
c = ?
```

```
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**step 3:** check all function alternatives

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swap (x,y) = (y,x)
```

```
a = (c,d)
```

```
b = (e,f)
```

```
c = ?
```

```
d = ?
```

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e = ?
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f = ?
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b = (d,c)
```

```
c = ?
```

```
d = ?
```

```
e = d
```

```
f = c
```

**step 3:** check all function alternatives

# List processing functions

- Students struggle to orchestrate list pattern-matching and list recursion, in particular when lists-of-lists are involved
- Effective remedy is to provide a design recipe (inspired by HtDP school)

# How to tackle recursion on lists?

1. Specify the name and write down a few cases:

```
my_function ... [] ... = ...
my_function ... [x1] ... = ...
my_function ... [x2, x1] ... = ...
my_function ... [x3, x2, x1] ... = ...
```

2. Specify the type of the function:

```
my_function :: ... [ ... ] ... -> ...
```

3. Make a case distinction based on the structure of your recursive data type.

- a) The list is empty (base case)

```
my_function ... [] ... = ...
```

depending on the problem, you may require more base cases

- b) The list is not empty (recursive case)

```
my_function ... [... : ...] ... = ...
```

depending on the problem, you may require more recursive cases



# How to tackle recursion on lists?

1. Specify the name and write down a few cases:

```
my_function ... [] ... = ...  
my_function ... [x1] ... = ...  
my_function ... [x2,x1] ... = ...  
my_function ... [x3,x2,x1] ... = ...
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my_function :: ... [ ... ] ... -> ...
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3. Make a case distinction based on the structure of your recursive data type.

- a) The list is empty (base case)

```
my_function ... [] ... = ...
```

- b) The list is not empty (recursive case)

```
my_function ... [... :xs] ... = r
```

where

```
r1 = my_function ... xs ...  
r = ... r1 ...
```

assume you have the result r1 for the smaller list xs

use r1 to solve problem for the argument list

# Length of a list

- Specify the name and write down a few cases:
  - length [] = 0
  - length [x1] = 1
  - length [x2, x1] = 2
  - length [x3, x2, x1] = 3

2. Specify the type of the function:

length :: [a] -> Int

3. Make a case distinction based on the structure of your recursive data type.

a) the list is empty

length [] = 0

b) the list is not empty

length [x : xs] = r

where

r1 = length xs

r = 1 + r1

} a  
bit  
verbose,  
transform

} length :: [a] -> Int  
length [] = 0  
length [\_ : xs] = 1 + length xs

use \_ to document that the head is not used in this alternative



# Initial segment of a list

- We wish to find the initial segment of a list (list without last element):

- `init [] = abort "init of []"`
- `init [x1] = []`
- `init [x2, x1] = [x2]`
- `init [x3, x2, x1] = [x3, x2]`
- ...

this function has 2 base cases

find the symmetry

# Initial segment of a list

- We wish to find the initial segment of a list (list without last element):

- `init [] = abort "init of []"`
- `init [x1] = []`
- `init [x2, x1] = [x2 : []]`
- `init [x3, x2, x1] = [x3 : [x2 : []]]`
- ...

this function has 2 base cases

find the symmetry

apply the steps...

```

init :: [a] -> [a]
init []      = abort "init of []"
init [_]    = []
init [x : xs] = [x : init xs]

```



# Higher-order functions

- Students find it hard to get a grip on higher-order functions
- Effective remedy is to show that a function 'is just another' parameter

*“It is also the goal for which functional programmers must strive - smaller and simpler and more general modules, glued together with the new glues we shall describe.”*

John Hughes, “Why Functional Programming Matters”, 1989

# Why smaller and simpler and more general functions matter...

- Smaller modules:
  - a module should concentrate on one piece of functionality
  - less prone to errors
  - easier to understand and reason about
- Simpler modules:
  - less prone to errors
  - easier to combine with other modules
- More general modules:
  - solve a recurring problem once and for all instead of ad-hoc
  - saves development time
  - errors rise from the application, not the implementation

## Functional programming:

- functions are compositional modules
- functions can be passed as parameter
- new functions can be computed

---

## Higher-order functions (HOF)

# Many functions look very similar

- `increase` `xs = [ x + 1` `\\ x <- xs ]`
- `decrease` `xs = [ x - 1` `\\ x <- xs ]`
- `squares` `xs = [ x * x` `\\ x <- xs ]`
- `avgs` `xs = [ (x + y) / (fromInt 2)` `\\ (x,y) <- xs ]`

```
inc x      = x + 1
dec x      = x - 1
square x   = x * x
avg (x,y)  = (x + y) / (fromInt 2)
```

# Many functions look very similar

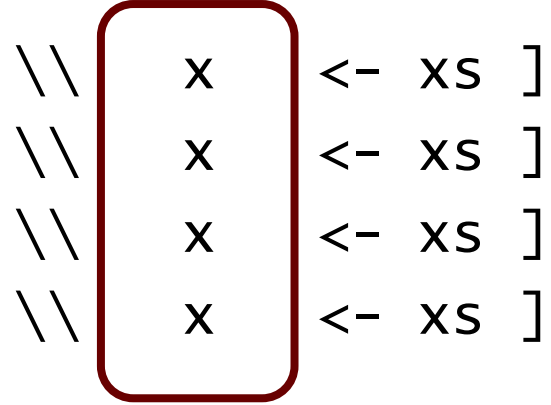
- `increase` `xs = [ inc x`
- `decrease` `xs = [ dec x`
- `squares` `xs = [ square x`
- `avgs` `xs = [ avg (x,y)`

Diagram illustrating the similarity in function signatures. A large rounded rectangle encloses the function names and their arguments: `inc x`, `dec x`, `square x`, and `avg (x,y)`. To the right, a vertical list of double backslashes `\\` separates this from another rounded rectangle containing the corresponding return types: `x`, `x`, `x`, and `(x,y)`. To the right of these return types is the assignment operator `<- xs ]`.

```
inc x      = x + 1
dec x      = x - 1
square x   = x * x
avg (x,y)  = (x + y) / (fromInt 2)
```

# Many functions look very similar

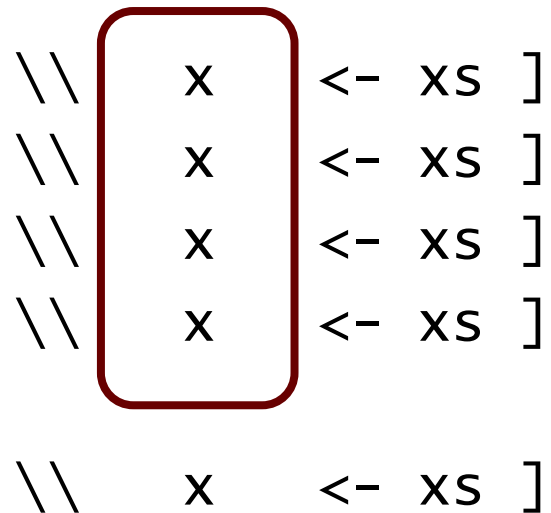
- `increase xs = [ inc x`
- `decrease xs = [ dec x`
- `squares xs = [ square x`
- `avgs xs = [ avg x`



```
inc x      = x + 1
dec x      = x - 1
square x   = x * x
avg (x,y)  = (x + y) / (fromInt 2)
```

# Many functions look very similar

- `increase xs = [ inc x`
  - `decrease xs = [ dec x`
  - `squares xs = [ square x`
  - `avgs xs = [ avg x`
- 
- `map f xs = [ f x`



```
inc x      = x + 1
dec x      = x - 1
square x   = x * x
avg (x,y)  = (x + y) / (fromInt 2)
```

# Many functions look very similar

- `increase xs = map inc xs`
- `decrease xs = map dec xs`
- `squares xs = map square xs`
- `avgs xs = map avg xs`
  
- `map f xs = [ f x | x <- xs ]`

```
inc x      = x + 1
dec x      = x - 1
square x   = x * x
avg (x,y)  = (x + y) / (fromInt 2)
```

- all applications of the same computational pattern: emphasize similarity
- ad-hoc part is identified: emphasize difference
  
- module with single purpose, and customizable due to HOF
  
- smaller and simpler functions
- easier to maintain
- can be reused in other context

# Many recursive functions look very similar

```

sum    :: [a] -> a | zero, + a
sum    []      = zero
sum    [x : xs] = x + sum xs

```

```

product :: [a] -> a | one, * a
product []      = one
product [x : xs] = x * product xs

```

```

and    :: [Bool] -> Bool
and    []      = True
and    [x : xs] = x && and xs

```

```

or     :: [Bool] -> Bool
or     []      = False
or     [x : xs] = x || or xs

```

```
sum    xs = foldr (+) zero xs
```

```
product xs = foldr (*) one xs
```

```
and    xs = foldr (&&) True xs
```

```
or     xs = foldr (||) False xs
```

```

foldr :: (a b -> b) b [a] -> b
foldr op e []      = e
foldr op e [x:xs] = op x (foldr op e xs)

```



# Higher-order functions

- Students find it hard to get a grip on higher-order functions
- Effective remedy is to show that a function ‘is just another’ parameter
- Using higher-order functions effectively is an acquired taste
  - in the ‘Segments: An alternative rainfall problem’<sup>1</sup> study, 32% of solutions use higher-order functions

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<sup>1</sup> ACHTEN, P. (2021). Segments: An alternative rainfall problem. *Journal of Functional Programming*, 31, E23. doi:10.1017/S0956796821000216

# Engaging projects

# SoccerFun<sup>1</sup>

- Implement a team of football players (or rather, their brain functions)
- Brain function:
  - :: (BrainInput,memory) -> (BrainOutput,memory)
- Brain input:
  - referee information
  - state of the ball
  - all other field players
  - player
- Brain output:
  - move, feint
  - kick ball, head ball
  - gain ball, catch ball, tackle



<sup>1</sup> Achten, P. (2008) Teaching functional programming with Soccer-Fun. In *Proceedings of the 1st Workshop on Functional and Declarative Programming in Education, FDPE'08*. Huch, F. & Parkin, A. (eds). Victoria, BC, Canada: ACM Press, pp. 61-72

# SoccerFun

- Can be tuned for full-fledged implementation to minimalistic approach
- Minimalistic approach can already be done after custom data types
- Recommended position is after lists / higher-order functions
- Students can concentrate on different aspects:
  - best written / structured code
  - best performing team
- Clever, unexpected solutions!

# General guidelines

# General guidelines

- Know the context and your students
- Construct a narrative for the course
- Anticipate hurdles and eliminate / lower them before they occur
- Don't forget yourself
  
- Thank you for your attention!