

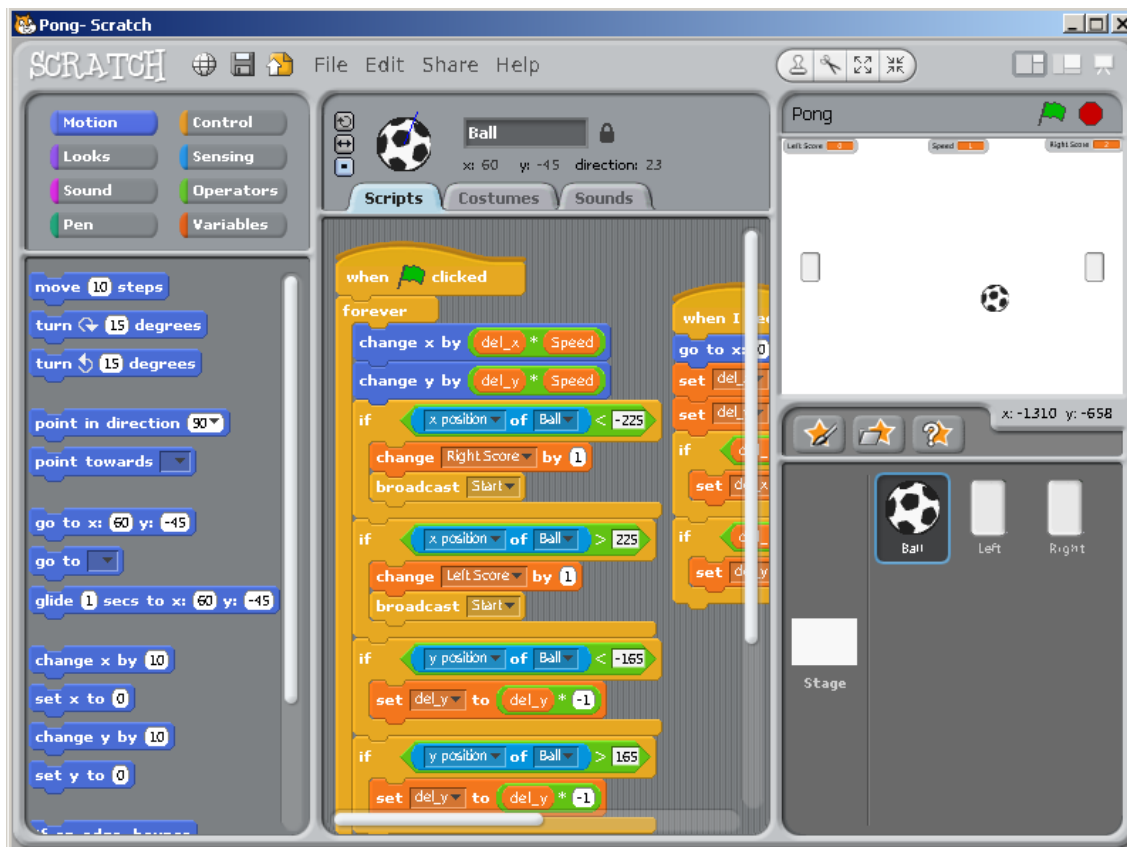
THEY ALREADY KNOW THE SYNTAX!

The Case for Spreadsheets in Programming Education

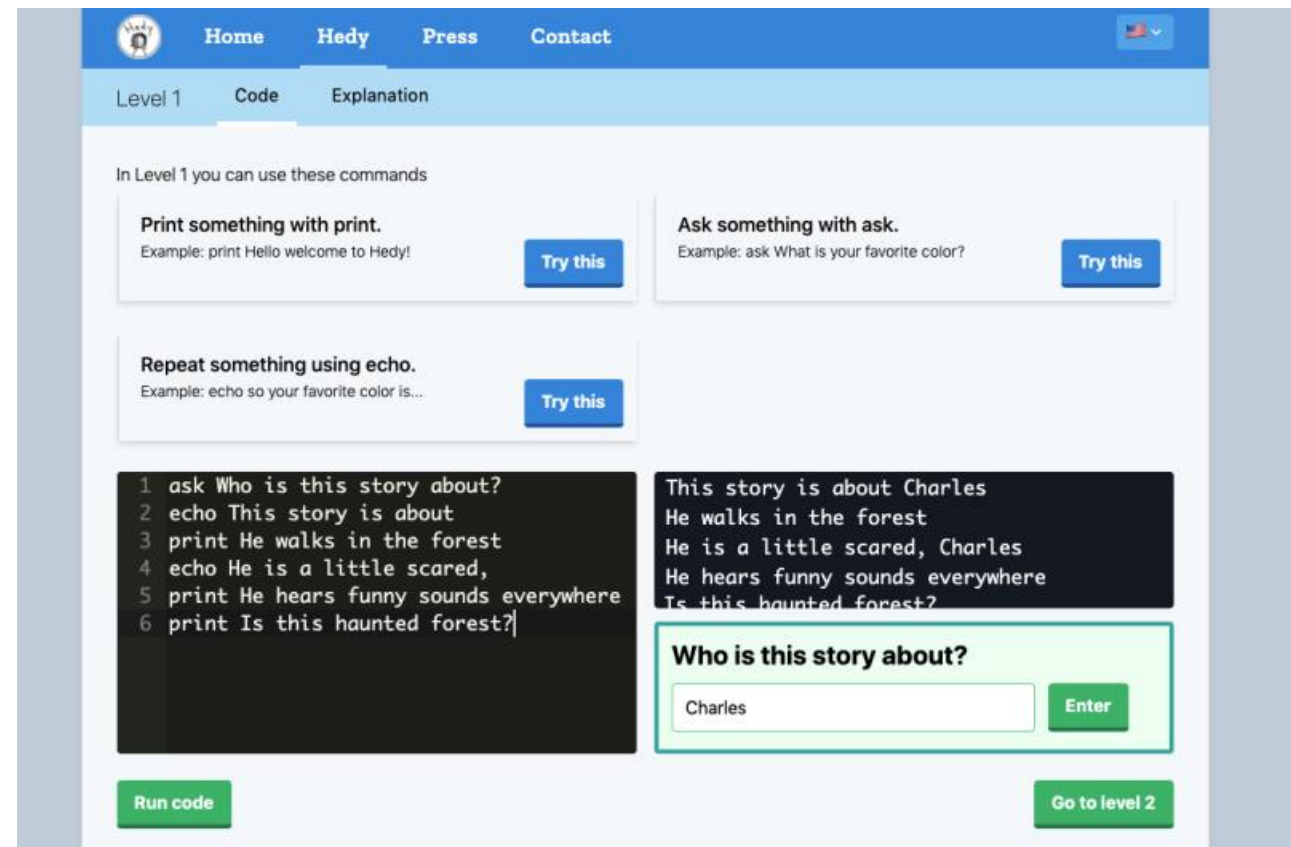
Learning Syntax is Known to be an Obstacle in Programming Education

Responses: **Block Coding** (replace syntax with shapes) & **Gradual Languages** (relaxed syntax rules)

Scratch [MIT Media Lab]

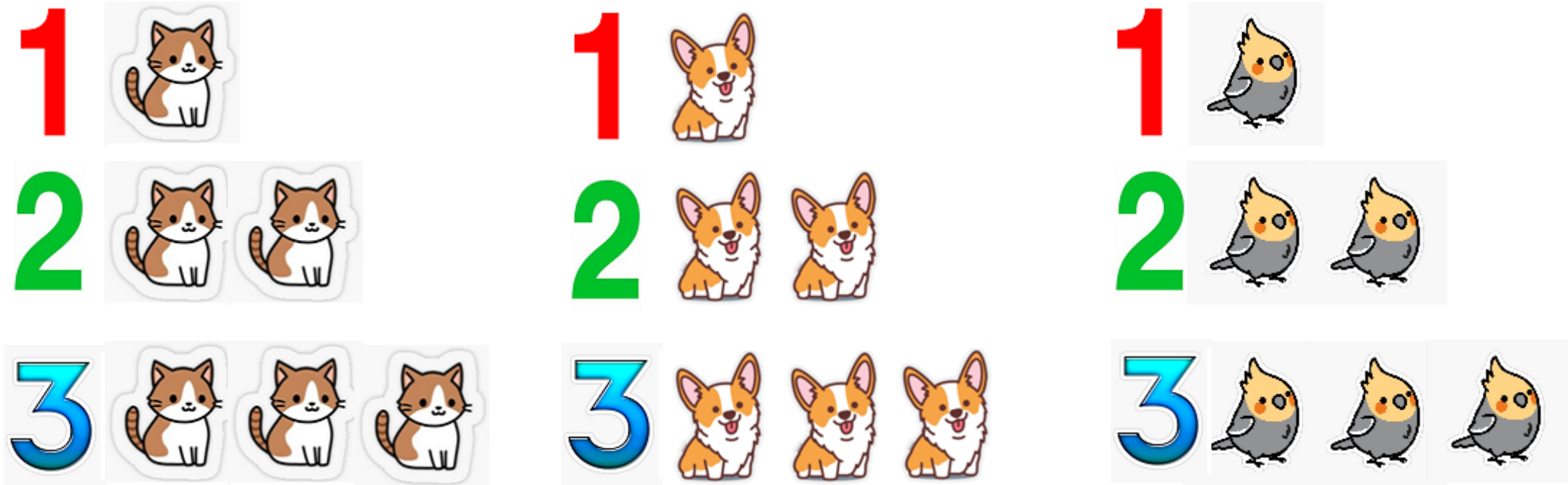


Hedy [Feliene Hermans, Leiden University]



Back to basics - let's reminisce our early computing education:

Math Abstractions ... Baby Steps



I am **one** baby syntax turtle →



Positional numeral system, operations, operator precedence (!) ...

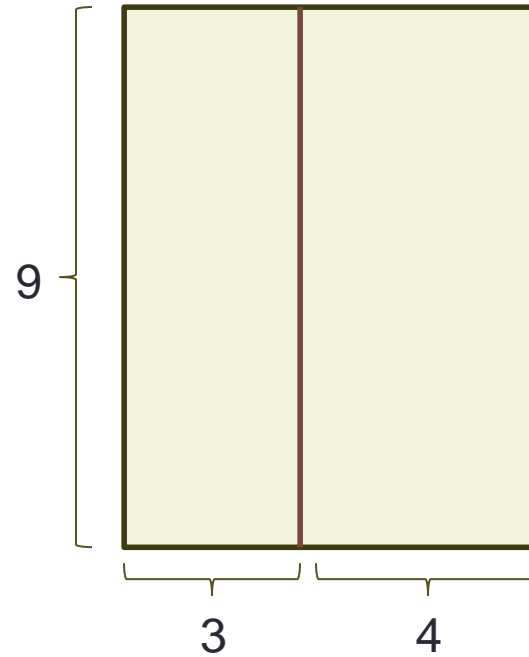
Math Abstractions ...

$$\begin{array}{r} 11 + \\ 31 \\ \hline 42 \end{array}$$

$$6 \times 7 = 42$$

$$2/7 + 3/7 = 5/7$$

$$1,000,000 - 997 = 999,003$$



$$9 \times (3 + 4) = 63$$



The turtles ... they keep ... multiplying (!) – here we are, at the end of middle school and early high school:

Math Abstractions

$$E = mc^2$$

$$y = m \cdot x + q$$

$$\sin^2 x + \cos^2 x = 1$$

$$\frac{\sin x}{\cos x} = \tan x$$

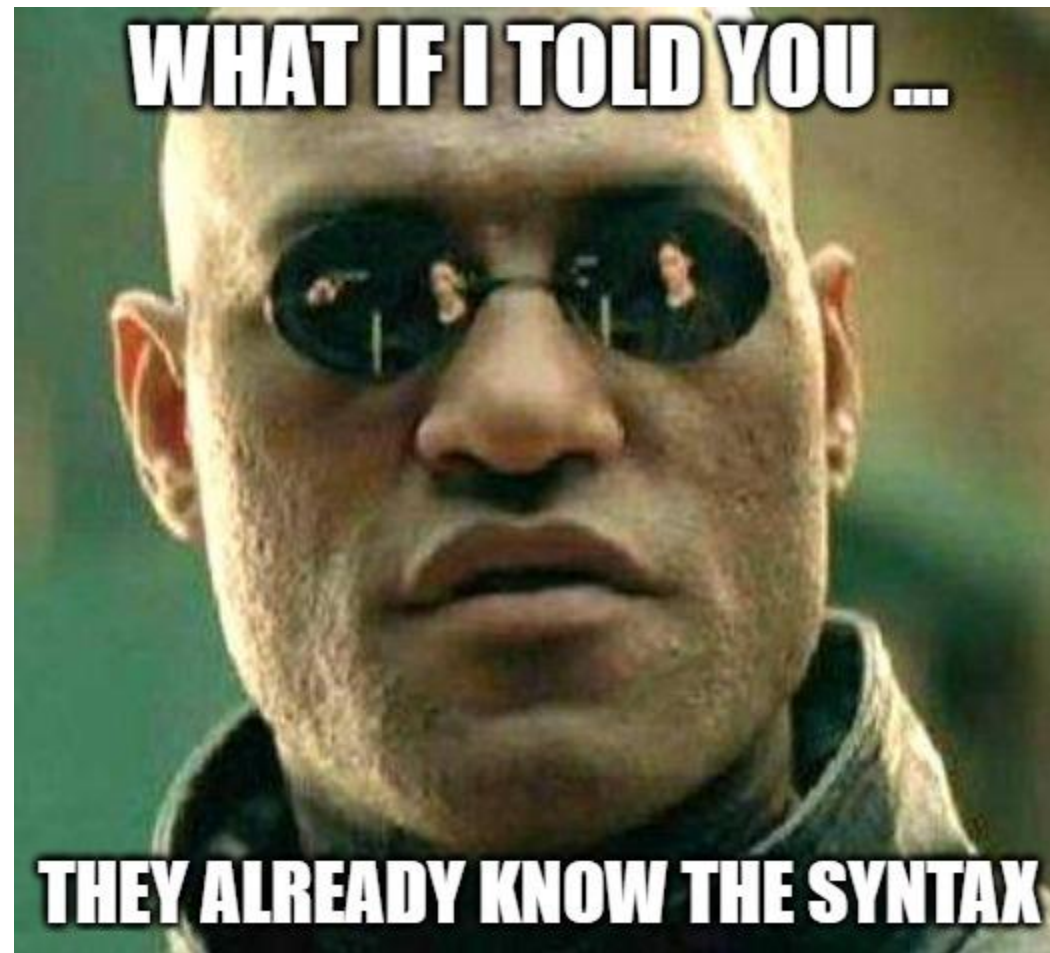
$$(a - b)(a + b) = a^2 - b^2$$

$$D = \sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2}$$

$$F = G \frac{m_1 m_2}{d^2}$$



How much *more* “syntax” must we feed students to start learning programming?



Note: G.A.M.N. stands for Generally accepted Math notation

Students Already Know (most of) The Syntax!

	Raise to Power	Multiplication	Division	Modulo	Addition	Subtraction	String Concatenation
MATH	x^y	$xy \mid x * y$	$x \div y \mid x / y$	$x \bmod y$	$x + y$	$x - y$	$xy \mid x \cdot y$
FORTRAN	<code>x ** y</code>	$x * y$	x / y	<code>mod(x, y)</code>	$x + y$	$x - y$	<code>x // y</code>
LISP	<code>(pow x y)</code>	<code>(* x y)</code>	<code>(/ x y)</code>	<code>(mod x y)</code>	<code>(+ x y)</code>	<code>(- x y)</code>	<code>(concatenate x y)</code>
C / C++	<code>pow(x, y)</code>	$x * y$	x / y	$x \% y$	$x + y$	$x - y$	$x + y$
Haskell	$x^y \mid x ** y$	$x * y$	x / y	<code>mod x y</code>	$x + y$	$x - y$	$x ++ y$
Python	<code>x ** y</code>	$x * y$	x / y	$x \% y$	$x + y$	$x - y$	$x + y$
Java	<code>Math.pow(x, y)</code>	$x * y$	x / y	$x \% y$	$x + y$	$x - y$	$x + y$
JavaScript	<code>x ** y</code>	$x * y$	x / y	$x \% y$	$x + y$	$x - y$	$x + y$
OCaml	<code>x ** y</code>	$x * y \mid x *. y$	$x / y \mid x /. y$	$x \bmod y$	$x + y \mid x +. y$	$x - y \mid x -. y$	x^y
MS-Excel	x^y	$x * y$	x / y	<code>mod(x, y)</code>	$x + y$	$x - y$	$x \& y$

∴ OK, sure, but those are *just expressions*. That's not programming, right? I mean, expressions aren't enough, right? ➔

A User-Centred Approach to Functions in Excel

30th June 2003

Simon Peyton Jones
Microsoft Research

Alan Blackwell
Cambridge University

Margaret Burnett
Oregon State University

“It may seem odd to describe a spreadsheet as a programming language. Indeed, one of the great merits of spreadsheets is that **users need not think of themselves as doing “programming”, let alone functional programming** — rather, they simply “write formulae” or “build a model”. However, one can imagine printing the cells of a spreadsheet in textual form, like this:

A1 = 3

A2 = A1-32

A3 = A2 * 5/9

and then it plainly is a (functional) program.”

“just expressions”: yes, to program with spreadsheets all the syntax you need to know is that of expressions!

But traditional spreadsheets have issues that make them unacceptable for education and other purposes ...

Critique of the Traditional Spreadsheet Core

- **Lack of functional abstraction**

- Considerable research work has been done on this
- December 3rd, 2020: Microsoft Research announced LAMBDA

- **Overly simplistic type system**

- All top-level variables must be a worksheet
- Worksheets are non-composable cell containers
- All cells are untype and must be referenced via coordinates
- **A1 notation should be considered harmful**

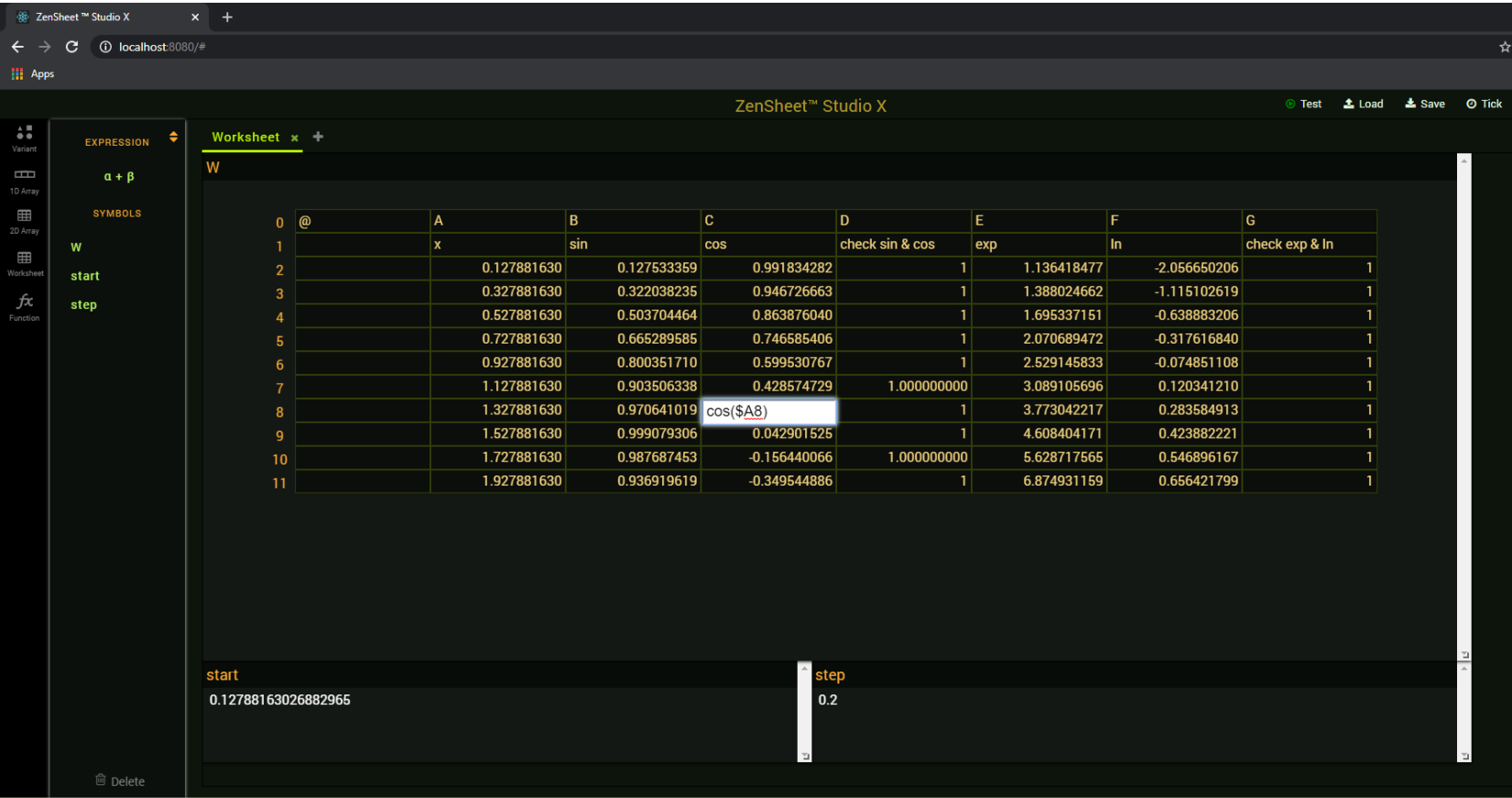
- **Entanglement of model and visualization**

- Worksheets are the only true variables of the core
 - They are containers that hold state, which includes unreduced expressions
- Worksheets are also the primary element of the presentation
 - They play an important role as UI layout managers

A language-centric redesign of spreadsheets has been shown to work well

ZenSheet / Lilly

ZenSheet supports composable data structures and functional abstraction.
2D arrays can be used as worksheets: it truly generalizes spreadsheets!



Types

- ZT.1) $T \rightarrow \text{null} \mid \text{error} \mid \text{bool} \mid \text{number} \mid \text{string}$
- ZT.2) $T \rightarrow \text{fun}(T, \dots, T) \Rightarrow T$
- ZT.3) $T \rightarrow \text{array}[\dots] \Rightarrow T$
- ZT.4) $T \rightarrow \text{struct}(T, \dots, T)$
- ZT.5) $T \rightarrow \text{lazy } T$
- ZT.6) $T \rightarrow \text{var}$
- ZT.7) $T \rightarrow \langle \text{symbol} \rangle$

Expressions

- XLS.1) $E \rightarrow ? \mid \langle \text{error} \rangle \mid \text{true} \mid \text{false} \mid \langle \text{number} \rangle \mid \langle \text{string} \rangle$
- XLS.3) $E \rightarrow \langle A1 \rangle \mid \langle \text{symbol} \rangle ! \langle A1 \rangle$
- ZSE.1) $E \rightarrow \langle \text{symbol} \rangle$
- ZSE.2) $E \rightarrow \lambda(T \langle \text{symbol} \rangle, \dots, T \langle \text{symbol} \rangle) \rightarrow E$
- ZSE.3) $E \rightarrow E(E, \dots, E)$
- ZSE.4) $E \rightarrow (E, \dots, E)$
- ZSE.5) $E \rightarrow [E, \dots, E]$
- ZSE.6) $E \rightarrow E[E, \dots, E]$
- ZSE.7) $E \rightarrow E:E$
- ZSE.8) $E \rightarrow E..E$
- ZSE.9) $E \rightarrow 'E'$

Actions

- ZSA.1) $A \rightarrow \text{type } \langle \text{symbol} \rangle = T;$
- ZSA.2) $A \rightarrow T \langle \text{symbol} \rangle := E;$
- ZSA.3) $A \rightarrow E := E;$

Listing 3: abstract syntax of Lilly

But to achieve this result we had to extend a traditional model of computing

Extending Christopher Strachey's Model

Traditional



- ◇ LVALUE
- ◇ CVALUE
- ◇ RVALUE

Assume φ has l-value, then:

- ◇ $RVALUE(\varphi) = CVALUE(LVALUE(\varphi))$

Reassignment:

- ◇ $\langle lhs \rangle := \langle rhs \rangle;$
 - ◇ Post: $CVALUE_p(\langle lhs \rangle) = RVALUE_b(\langle rhs \rangle)$

Lilly

- ◇ LVALUE
- ◇ CVALUE
- ◇ RVALUE

Assume φ has l-value, then:

- ◇ $RVALUE(\varphi) = RVALUE(CVALUE(LVALUE(\varphi)))$

Reassignment:

- ◇ $\langle lhs \rangle := \langle rhs \rangle;$
 - ◇ Post: $CVALUE_p(\langle lhs \rangle) = RVALUE_b(\langle rhs \rangle)$

Lilly natively supports a function *formula*(φ) that returns $CVALUE(\varphi)$ without computing $RVALUE(\varphi)$

Experience Report: Lilly in Action

Our Experience Report: Complementing Programming Education with Lilly and ZenSheet

- Setting
 - Online classes (university is still closed) with no technical support
 - Students set up their own lab, with material and assistance provided by yours truly
 - GitHub (<https://github.com/>), MSYS2 (<https://www.msys2.org/>), ...
 - Students must learn C in 10 weeks
 - Nearly all of them have no programming experience
 - An old professor and I have been advocating to modernize the curriculum
 - ... in fact, we have already started to do so, under the RADAR
 - ... and have been gaining support from researchers and institutions
 - Lilly will officially be part of the course in the period that starts next week



Lilly REPL Session

The response to an **action** can be an ACK, followed by an echo of the command, or an ERROR with a descriptive message.

The response to an **expression** can be an OK followed by a description of the reduction, or an ERROR with a descriptive message.

Variable `z` in this session is initialized with a quoted expression, therefore inferred to have a lazy type

ZenSheet REPL - Beta 0.1

Connecting via net protocol to localhost:3899 ...

< ZR > :: data := [0, 1, 42, 67, 3, 7, 997, 8];

ACK: :: data := [0, 1, 42, 67, 3, 7, 997, 8];

< ZR > filter(fn(x) -> x < 50, data)

OK: filter(fn(x) -> x < 50, data) ==> [0, 1, 42, 3, 7, 8]

< ZR > :: predicate := fn(x) -> x < 50;

ACK: :: predicate := fn(x) -> x < 50;

< ZR > :: z := 'filter(predicate, data)';

ACK: :: z := 'filter(predicate, data)';

< ZR > z

OK: z ==> [0, 1, 42, 3, 7, 8]

< ZR > predicate := fn(x) -> x % 2 = 0;

ACK: predicate := fn(x) -> x % 2 = 0;

<13: 7/7 ==> 7> z

OK: z ==> [0, 42, 8]

Higher Order Functional Abstraction Example

User implementation of `map`, `filter`, `fold`

```
// map
:: mapz := fn(f, seq) ->
    if(empty(seq), seq, cons(f(head(seq)), mapz(f, tail(seq))));

// filter
:: filterz := fn(pred, seq) ->
    if(empty(seq), seq, if(pred(head(seq)), cons(head(seq), filterz(pred, tail(seq))), filterz(pred, tail(seq))));

// fold
:: foldz := fn(mfn, init, seq) ->
    if(empty(seq), init, foldz(mfn, mfn(init, head(seq)), tail(seq)));
```

Higher Order Functions: a Reactive Pipeline Model

```
///
/// hof.sym
///
/// higher order functions example: reactive pipeline model
///

array[] => array[] => lazy double input := [
  ['uniform()', 'uniform()', 'uniform()', 'uniform()', 'uniform()', 'uniform()'],
  ['uniform()', 'uniform()', 'uniform()', 'uniform()', 'uniform()', 'uniform()'],
  ['uniform()', 'uniform()', 'uniform()', 'uniform()', 'uniform()', 'uniform()']
];

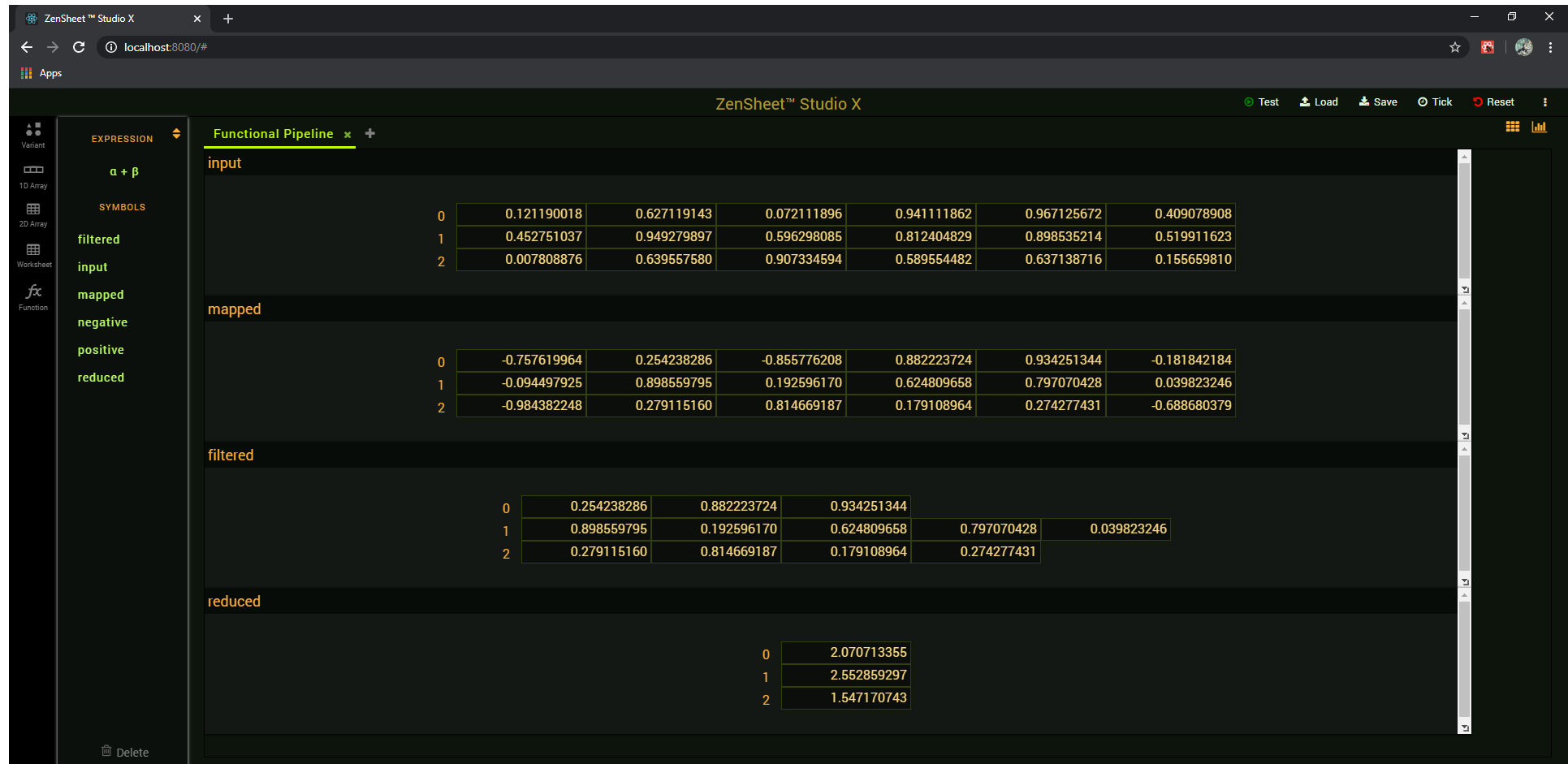
lazy var negative := './.(x) -> x < 0';
lazy var positive := './.(x) -> x >= 0';

lazy var mapped := 'map(/.(row) -> map(/.(x) -> 2 * x - 1, row), input)';

lazy var filtered := 'map(/.(row) -> filter(positive, row), mapped)';

lazy var reduced := 'map(/.(row) -> sum(row), filtered)';
```

Reactive Pipeline Model – rendered in ZenSheet UI



sys.cycle(0);

Findings and Preliminary Decisions

- Post grade survey: nearly all students reported that Lilly was valuable or very valuable to their learning experience. No one considered it detrimental.
- **Tech issues: deployment turned out to be even more of a challenge than anticipated 😞**
- Focusing on concepts and paradigms, showing how they are supported in different languages, reduces the “are we learning the right language” worries.
- Adding **Lilly** and **JavaScript** appears to help overcome syntax-related anxiety.
- We are replacing C with **C++** (already used C++ last trimester) even more.
- We also plan to use **Haskell** to show examples of parametric polymorphism.

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